

Security

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The Security Problem

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 - How do we provide **controlled access** to programs and data **stored in a computer system**?

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 - How do we provide **controlled access** to programs and data **stored in a computer system**?

- ▶ **Security** requires both **protection system** and the consideration of the **external environment** within which the system operates.

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- ▶ **Threat** is potential **security violation**.
- ▶ **Attack** is attempt to **breach security**: **accidental** or **malicious**.

Security Violation Categories

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- ▶ **Theft of service**: unauthorized **use** of resources
- ▶ **Denial of Service (DoS)**: prevention of **legitimate use**

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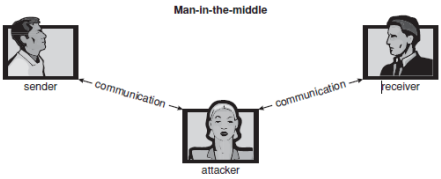
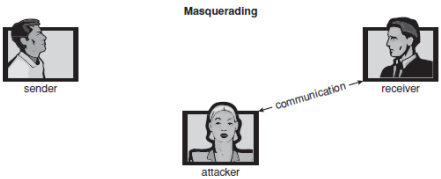
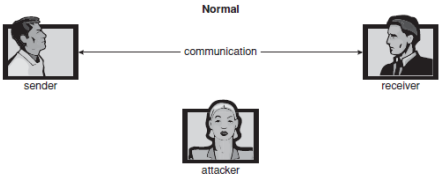
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- ▶ **Session hijacking**: intercept an **already-established** session to **bypass authentication**

Standard Security Attacks



- ▶ **Security** must occur at **four levels** to be effective:
 - **Physical**: data centers, servers, connected terminals
 - **Human**: only appropriate users have access to the system
 - **OS**: protection mechanisms, debugging
 - **Network**: intercepted communications, interruption, DOS

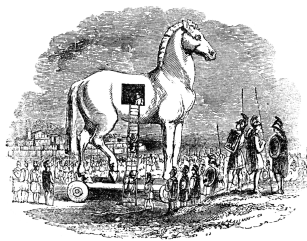
Program Threats

- ▶ Many variations, many names
 - Trojan horse
 - Trap door
 - Logic bomb
 - Stack and buffer flow
 - Viruses

Trojan Horse

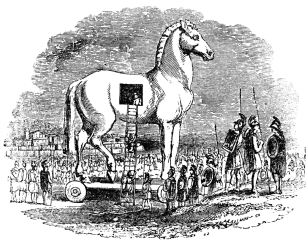
Trojan Horse (1/2)

- ▶ Code segment that **misuses** its **environment**.



Trojan Horse (1/2)

- ▶ Code segment that **misuses** its **environment**.
- ▶ Exploits mechanisms for **allowing programs written by users** to be **executed by other users**.
- ▶ Example:
 - A **text-editor** program has a code to **search the file** to be edited for certain keywords.
 - If any are found, the **entire file** may be copied to a special area accessible to the creator of the text editor.



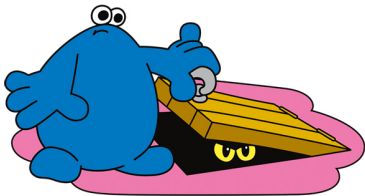
Trojan Horse (2/2)

- ▶ Variation of Trojan horse:
- ▶ Emulating a login program
- ▶ **Spyware**: accompanies a program that the user has installed.
 - Download ads to display on the user's system
 - Create pop-up browser windows when certain sites are visited
 - Capture information from the user's system and return it to a central site: covert channel

Trap Door

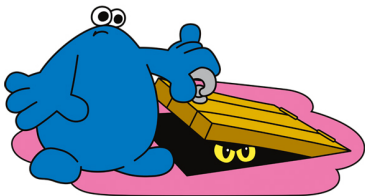
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Trap Door

- ▶ A **hole** in the software that **only designer** of a program is capable of using.
- ▶ Could be included in a **compiler**.
- ▶ Difficult to detect: we have to **analyze all the source code** for all components of a system.



Logic Bomb

Logic Bomb

- ▶ Program that **initiates a security incident** under **certain circumstances**.
- ▶ Hard to detect: because under **normal operations**, there would be **no security hole**.



Stack and Buffer Flow

Stack and Buffer Overflow (1/6)

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Stack and Buffer Overflow (1/6)

- ▶ The **most common way** for an **attacker** to **gain unauthorized access** to the target system.
- ▶ The attack **exploits a bug** in a program.
 - E.g., The programmer neglected to code **bounds checking on an input field**.
 - The attacker sends **more data** than the program was expecting.
 - The attacker can write a program to do the next page steps.

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Stack and Buffer Overflow (2/6)

- ① **Overflow an input field**, for example, a **web-page form** expects a **user name**, until it writes into the **stack**.
- ② **Overwrite** the current **return address** on the stack with the address of the **exploit code** loaded in step 3.
- ③ Write a simple **set of code** for the next space in the stack that includes the commands that the **attacker wishes to execute**, for instance, spawn a shell.

Stack and Buffer Overflow (3/6)

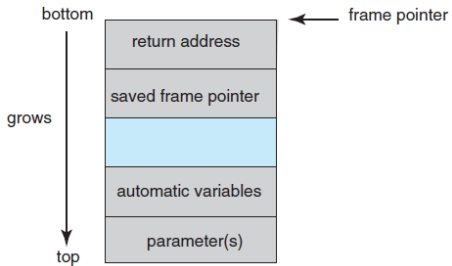
```
#include <stdio.h>
#define BUFFER_SIZE 256

int main(int argc, char *argv[])
{
    char buffer[BUFFER_SIZE];

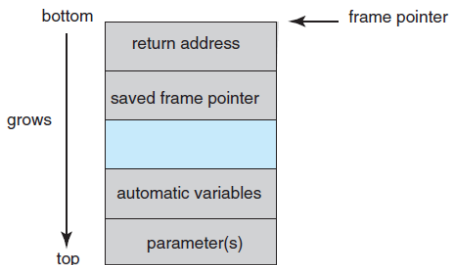
    if (argc < 2)
        return -1;
    else {
        strcpy(buffer,argv[1]);
        return 0;
    }
}
```

- ▶ Lack of bounds checking
- ▶ If the command line input is longer than `BUFFER_SIZE`?

Stack and Buffer Overflow (4/6)

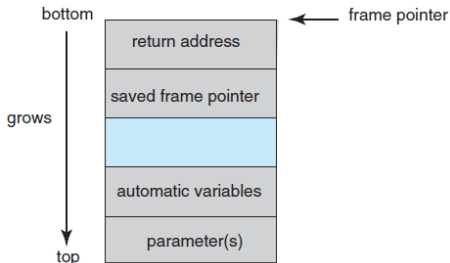


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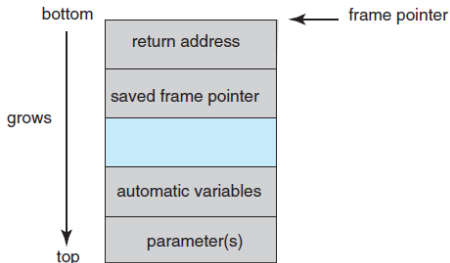
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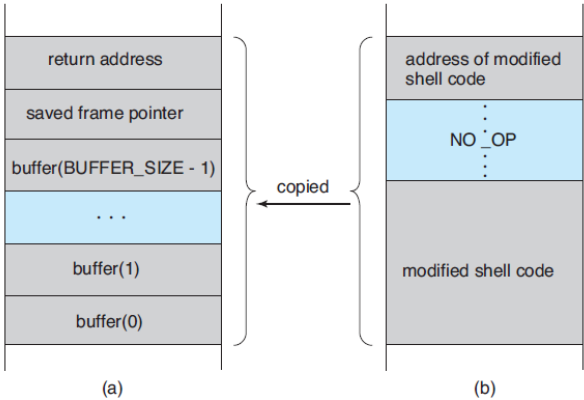
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- ▶ **Automatic variables:** variables defined locally to the function.
- ▶ **Frame pointer:** the address of the beginning of the stack frame.
 - Can vary during the function call.
- ▶ **Return address:** where to return control once the function exits.

Stack and Buffer Overflow (5/6)



Stack and Buffer Overflow (6/6)

```
#include <stdio.h>

int main(int argc, char *argv[])
{
    execvp(\"\bin\sh\", \"\bin \sh\", NULL);
    return 0;
}
```

- ▶ The cracker could **replace the return address** with the **address of the code segment** containing the **attacking program**.

Viruses

- ▶ Code fragment **embedded** in legitimate program.
- ▶ **Self-replicating**: designed to **infect other computers**.
- ▶ Very **specific** to CPU architecture, OS, applications.
- ▶ Usually borne via **email or as a macro**.

System and Network Threats

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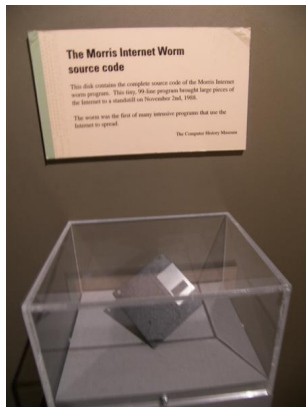
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- ▶ They create a situation in which the **OS resources and user files** are **misused**.
- ▶ They include:
 - **Worms**
 - **Port scanning**
 - **Denial of Service (DoS)**

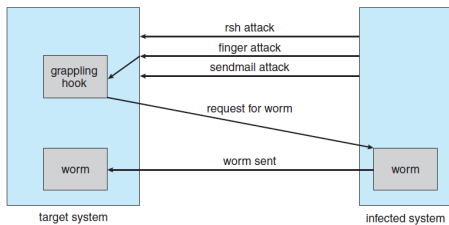
Worms (1/2)

- ▶ An **standalone program** that **replicates itself** in order to **spread** to other computers.
- ▶ Often, it uses a **computer network** to spread itself.
- ▶ **Unlike a computer virus**, it does **not need to attach itself** to an existing program.
- ▶ The **Morris worm** (the **Internet worm**) is the **first computer worms** distributed via the Internet: 1988



Worms (2/2)

- ▶ **Morris worm**: exploited UNIX networking features in `rsh` and bugs in `finger` and `sendmail` programs.
- ▶ It made up of **two programs**: a `grappling hook` program and the `main program`.
- ▶ The `grappling hook` program **uploaded main worm program**.



Port Scanning

- ▶ A means to **detect a system's vulnerabilities** to attack.
- ▶ Automated attempt to **connect to a range of ports** on IP addresses.
- ▶ Detection of **answering service protocols**, **OS** and **version** running on system.

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- ▶ Detection of **answering service protocols**, **OS** and **version** running on system.
- ▶ **nmap scans all ports** in a given IP range for a response
- ▶ **nessus** has a **database of protocols and bugs** (and exploits) to apply against a system.

Denial of Service (DoS)

- ▶ **Overload** the targeted computer **preventing** it from doing any useful work.
- ▶ **Distributed denial-of-service (DDoS)** come from **multiple sites at once**.
- ▶ Consider **traffic to a web site**.
- ▶ CS students writing bad **fork()** code.

Cryptography

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- ▶ Similarly, a **sender** can **encode its message** so that only a computer with a **certain key** can decode the message.

Encryption (1/2)

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 - For each $k \in K$, E_k is a function for generating ciphertexts from messages.

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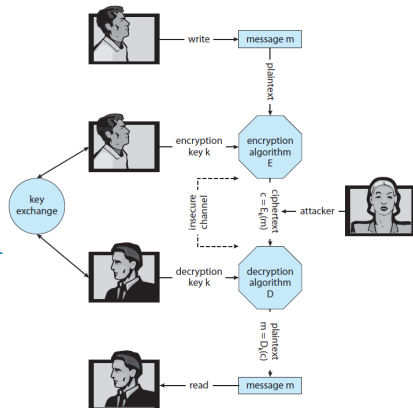
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 - ▶ A function $D : K \rightarrow (C \rightarrow M)$
 - For each $k \in K$, D_k is a function for generating messages from ciphertexts.

Encryption (2/2)

- ▶ An encryption algorithm **must provide** this **essential property**:
 - Given a ciphertext $c \in C$, a computer can compute m such that $E_k(m) = c$ only if it possesses k .
- ▶ Thus, a computer holding k can **decrypt ciphertexts** to the **plaintexts**.

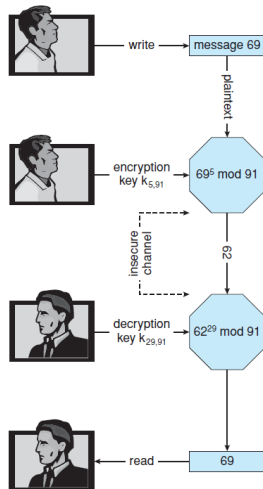
Symmetric Encryption

- ▶ Same key used to encrypt and decrypt.
 - k must be kept secret.
- ▶ E.g., DES, Triple-DES, AES, RC4



Asymmetric Encryption

- ▶ Each user having **two keys**:
 - **Public key**: published key used to **encrypt** data.
 - **Private key**: key known only to individual user used to **decrypt** data.
- ▶ Most common is **RSA** block cipher.



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 - A function $S : K \rightarrow (M \rightarrow A)$: for each $k \in K$, S_k is a function for **generating authenticators from messages**.
 - A function $V : K \rightarrow (M \times A \rightarrow \{\mathbf{true}, \mathbf{false}\})$: for each $k \in K$, V_k is a function for **verifying authenticators on messages**.

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- ▶ Thus, computer holding k can generate authenticators on messages so that any other computer possessing k can verify them.
- ▶ Practically, if $V_k(m, a) = \text{true}$ then we know m has not been modified and that send of message has k .

Authentication - Hash Function (1/2)

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- ▶ E.g., **MD5** (128-bit hash), and **SHA-1** (160-bit hash)

Authentication - Hash Function (2/2)

- ▶ Not useful as authenticators.
 - For example $H(m)$ can be sent with a message.
 - But if H is known someone could modify m to m' and recompute $H(m')$ and modification not detected
 - So must authenticate $H(m)$: MAC and digital signature

Authentication - MAC

- ▶ Symmetric encryption used in message-authentication code (MAC) authentication algorithm.
- ▶ Cryptographic checksum generated from message using secret key.

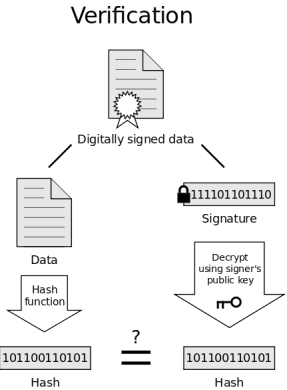
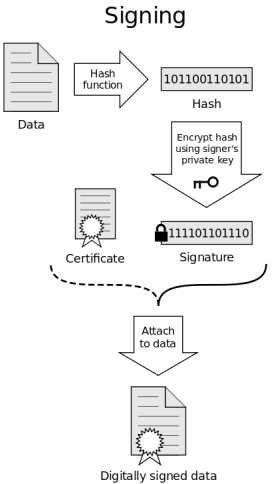
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- ▶ Note that k is needed to compute both S_k and V_k , so anyone able to compute one can compute the other end digital signature.

Authentication - Digital Signature (1/2)

- ▶ Based on **asymmetric keys** and digital signature algorithm.
- ▶ **Authenticators** produced are **digital signatures**.
- ▶ **Anyone** can **verify** authenticity of a message.
- ▶ In a digital-signature algorithm, computationally **infeasible** to derive k_s from k_v .
- ▶ k_v is the **public key** and k_s is the **private key**.

Authentication - Digital Signature (2/2)



If the hashes are equal, the signature is valid.

► Symmetric key

- Sometimes done **out-of-band**, e.g., via a paper document or a conversation
- If a user wants to communicate with N other users privately: she needs N keys.

Authentication - Key Distribution

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- A user needs only **one private key**, no matter how many other people she wants to communicate with.
- **Man-in-a-middle** attack.
- What we need is a **proof of who owns a public key?** **digital certificate**

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Authentication - Digital Certificates

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 - The **digital signature** of an entity that has **verified** the certificate's contents are **correct**.

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- ▶ Certificate authority are **trusted party** - **their public keys** included with web **browser distributions**.

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 - The user's **possession of something** (a key or card)
 - The user's **knowledge of something** (a user identifier and password)
 - An **attribute of the user** (fingerprint, retina pattern, or signature)

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 - Use algorithm **easy to compute** but **difficult to invert**.
 - Only **encrypted password stored**, never decrypted.
- ▶ **One-time** passwords
 - Use a function **based on a seed** to compute a password, both user and computer.

Implementing Security Defenses

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- ▶ **Security policy** describes **what is being secured**.
- ▶ **Vulnerability assessment** compares **real state** of a system/network to the **security policy**.
- ▶ **Intrusion detection systems** to detect **attempted or successful intrusions**.
 - **Signature-based detection** spots **known bad patterns**.
 - **Anomaly detection** spots **differences from normal behavior**.

▶ Virus protection

- Searching all programs or programs at execution for known virus patterns.

Implementing Security Defenses (2/2)

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 - Searching all programs or programs at execution for known virus patterns.

- ▶ Auditing, accounting, and logging of all or specific system or network activities.

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- ▶ User authentication: password

Questions?