

CLive: Cloud-assisted P2P Live Streaming

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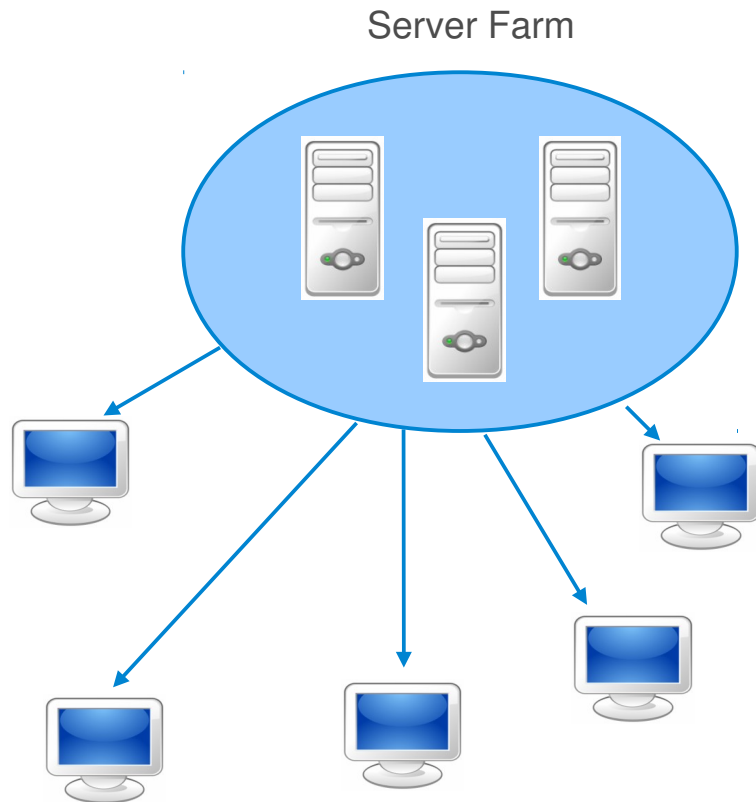
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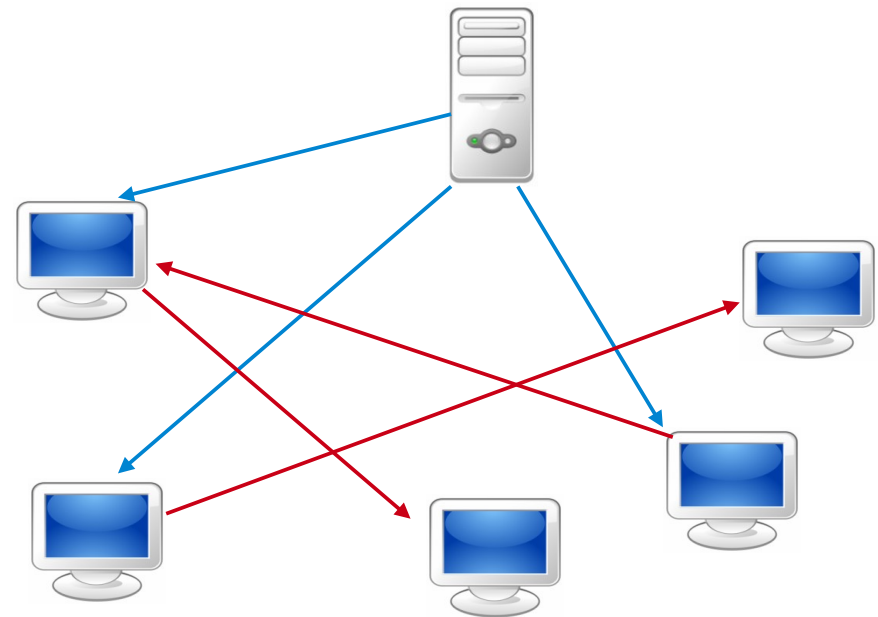
The Problem

Existing P2P Streaming Solutions

Client-Server



Peer-to-Peer



Problem?

- Bottlenecks in P2P video streaming systems: upload bandwidth
- A potential solution: P2P network is assisted by a cloud computing.

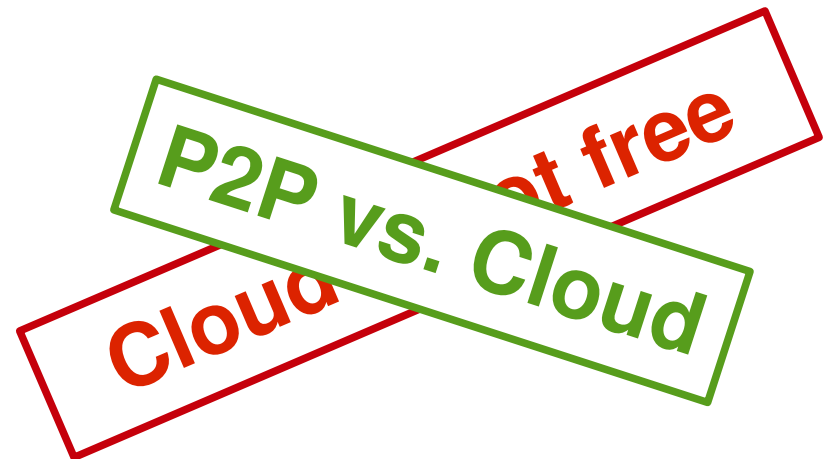
Problem?

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- A potential solution: P2P network is assisted by a cloud computing.

Cloud is not free

Problem?

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- A potential solution: P2P network is assisted by a cloud computing.



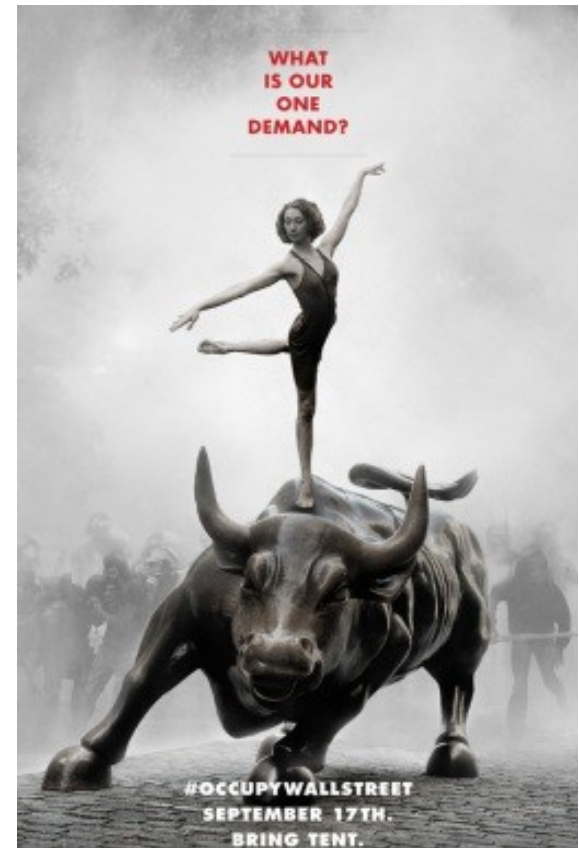
P2P vs. Cloud

- P2P pros and cons
 - P2P resources are cheap
 - Churn may compromise availability
- Cloud pros and cons
 - Superior availability
 - Cloud resources are not free



If You Cannot Beat Them, ...? ;)

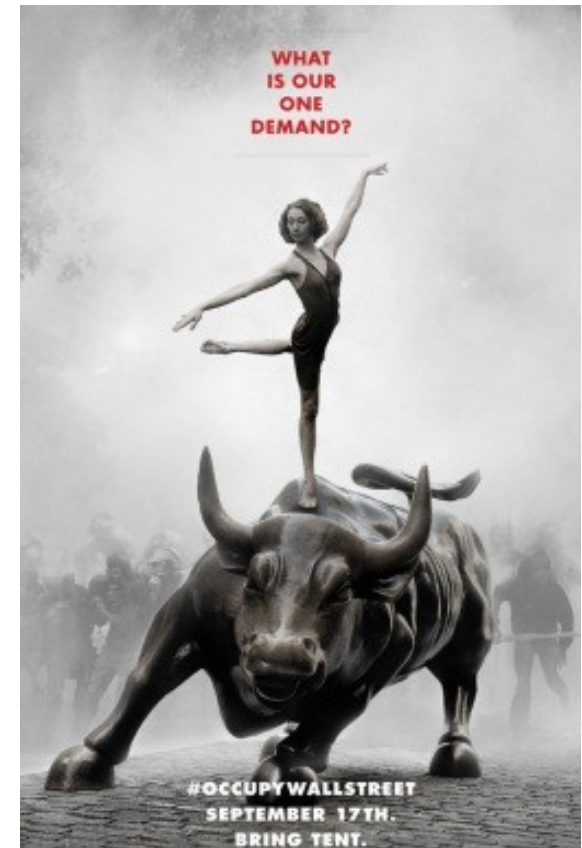
- The **cloud** as a **support group** for **P2P**.
- **Reduce** the number of (costly) **cloud interactions** as much as possible.



If You Cannot Beat Them, ...? ;)

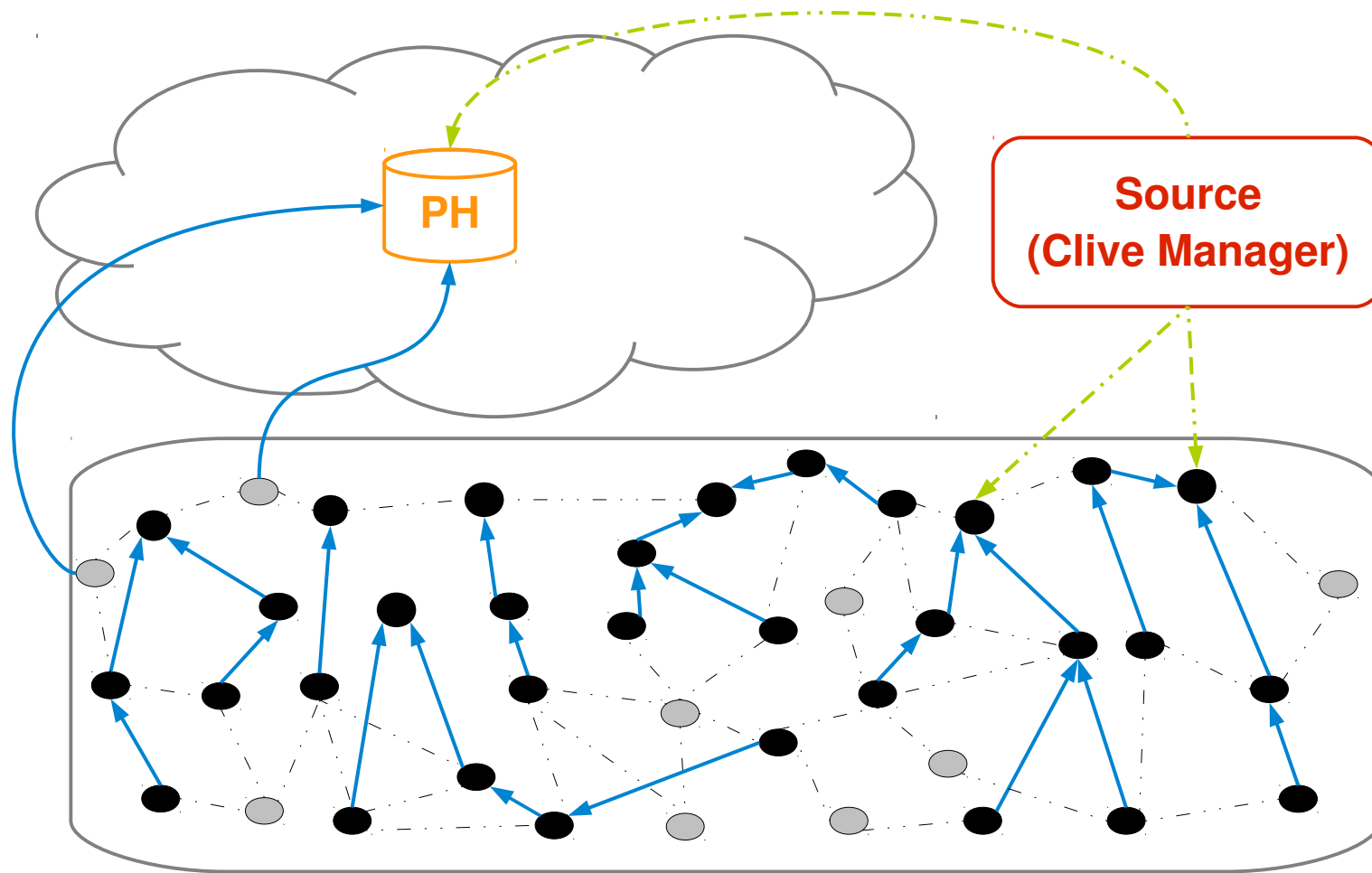
- The cloud as a support group for P2P.
- Reduce the number of (costly) cloud interactions as much as possible.

The problem to be solved becomes **minimizing the economical cost**, provided that a **set of constraints on QoS is satisfied**.

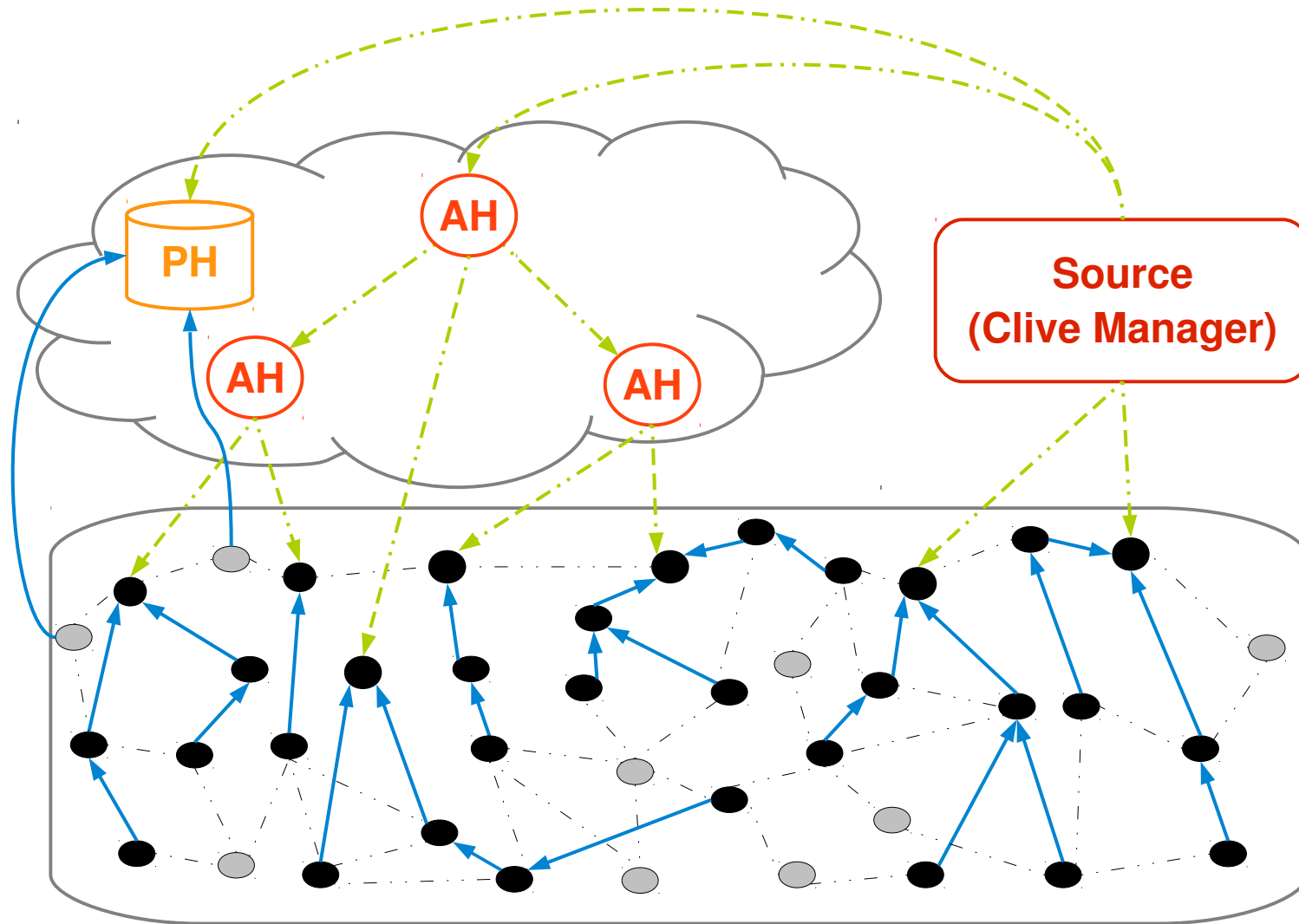


CLive

Baseline Model

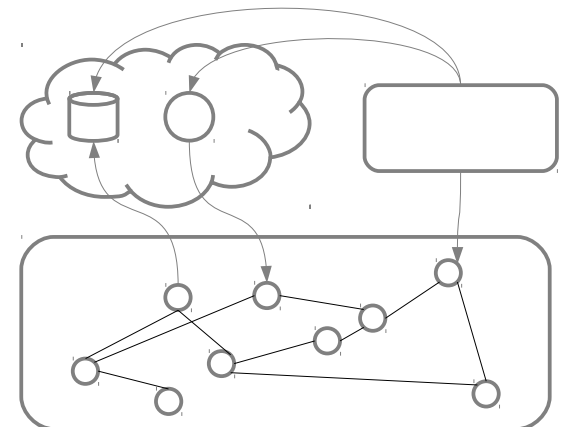


Enhanced Model



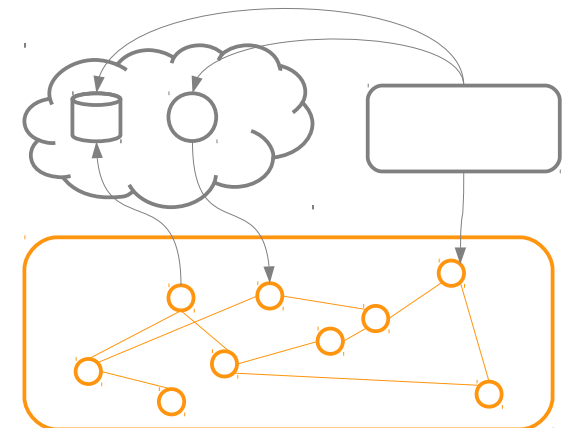
CLive Components

- P2P streaming overlay (swarm)
- Media source
- Passive Helpers
- Active Helpers
- Management component



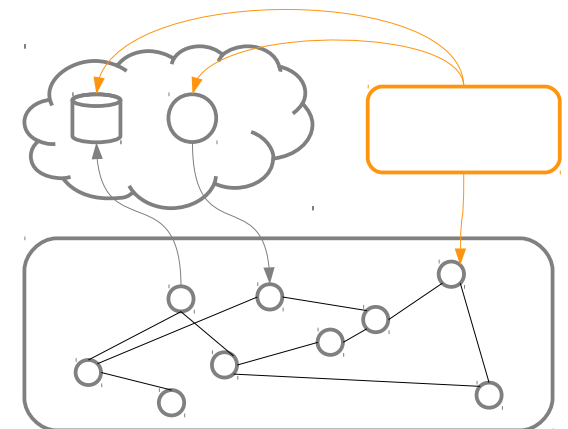
CLive Components: P2P Streaming Component

- We assume that nodes in the system use a **mesh-pull** model for data distribution.
- Nodes periodically send their **buffer maps** to their neighbours.
- The other nodes **pull** the required chunks from those nodes who own the chunks.



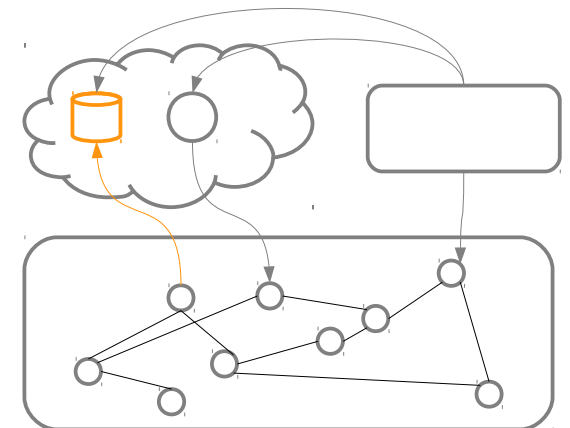
CLive Components: Media Source

- A **media source** is a node that generates data chunks and **pushes** them to the swarm.



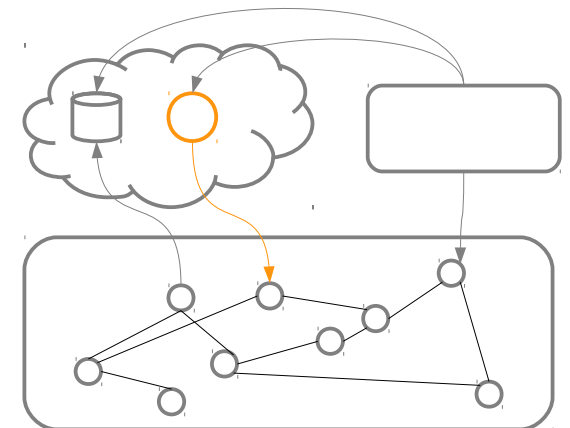
CLive Components: Passive Helper

- PH is a passive element that plays the role of a **data storage**, e.g., **Amazon S3**.
- The source **pushes chunks to PH**, as they are generated.
- The swarm nodes **pull the missed chunks** from it.
- We assume that a PH can serve as many requests as it receives.



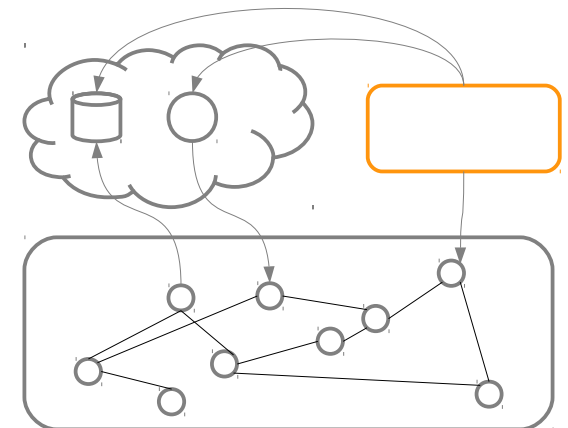
CLive Components: Active Helper

- AHs are active elements, e.g., Amazon EC2, that cooperate with other swarm nodes to accelerate the data dissemination.
- The source pushes chunks to AH, as they are generated.
- AHs forwards chunks to the swarm and other AH, as they are received.



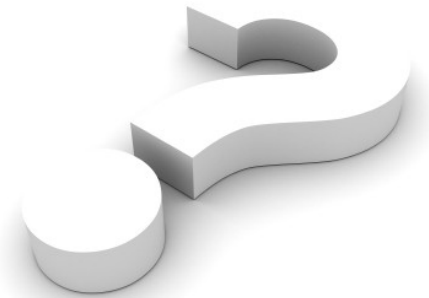
CLive Components: Management

- Participates in a **gossip algorithm** to **estimate the available resources** in the system.
- **Adds/removes AHs** to/from the system, based on the estimation.



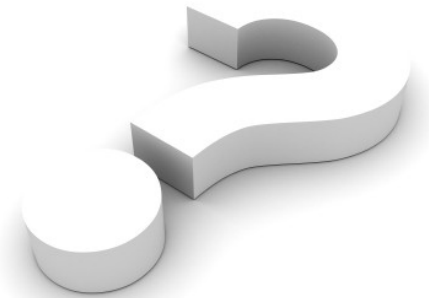
Two Main Questions in CLive?

- How to estimate the extra load in the overlay?
- How to relay the load to cloud with a minimum cost?



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Main Idea

- **Infected nodes**: the number of nodes that can be served with the **existing resources** in the system.
- $\text{load} = \text{swarm size} - \text{infected nodes}$

The Swarm Size

- The swarm size estimation is easy: gossip-based aggregation

The Infected Nodes

- It is shown that each streamed chunk through mesh overlays follows a **tree-based diffusion pattern** [1].
- To compute the number of **infected nodes**, **we model a diffusion tree for a chunk** and estimate the number of nodes in that tree.

[1] B. Biskupski, M. Schiely, P. Felber, and R. Meier, "Tree-based analysis of mesh overlays for peer-to-peer streaming," in Proc. of the 8th int. conf. on distr. app. and interoperable systems. Springer, 2008.

Estimate the Number of Nodes in a Tree

- Estimate the tree depth
- Estimate the fan-out distribution

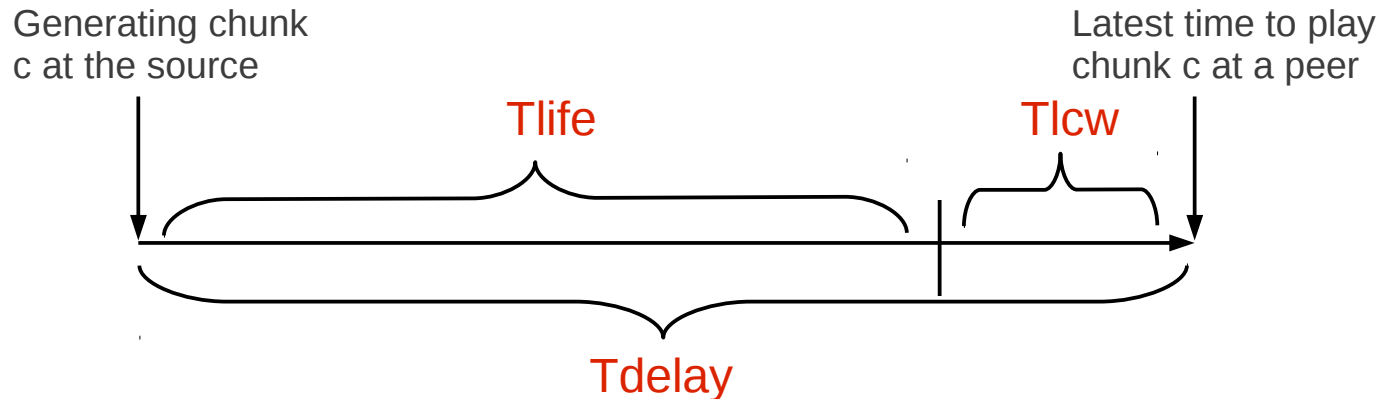
Estimate the Tree Depth (1/2)

- **Last Chance Window (LCW)**: A predefined number of data chunks that must be buffered ahead of the playback point.
- The missed chunks are **downloaded directly from the PH**, when the number of data chunks that are buffered ahead is **less than LCW**.
- However, if the **LCW buffer is full**, nodes pull the required chunks from the **swarm**.

Estimate the Tree Depth (2/2)

- T_{delay} : maximum acceptable latency.

- $T_{\text{life}} = T_{\text{delay}} - T_{\text{lcw}}$



- If a node can not receive a chunk in T_{life} , it pulls that chunk from PH.

- The tree depth

- $D_{\text{max}} = T_{\text{life}} / T_d$
- T_d : average latency among the peers

Estimate the Fan-out Distribution

- The upload bandwidth distribution
 - Adam2
 - Gossip-based aggregation

Estimate the Number of Nodes in a Tree

- Estimate the tree depth
- Estimate the fan-out distribution



n_{tree}

The Number of Infected Nodes

- $N_{inf} = (\text{the number of trees}) \cdot n_{tree}$
- The number of trees: source fanout + AHs fanout

Two Main Questions in CLive?

- How to estimate the extra load in the overlay?
- How to relay the load to cloud with a minimum cost?



AH Cost

- AH cost in one round: $C_{ah} = C_{vm} + m \cdot C_{chunk}$
 - C_{vm} : virtual machine cost
 - C_{chunk} : chunk transfer cost
 - m : number of chunks that one AH upload per round

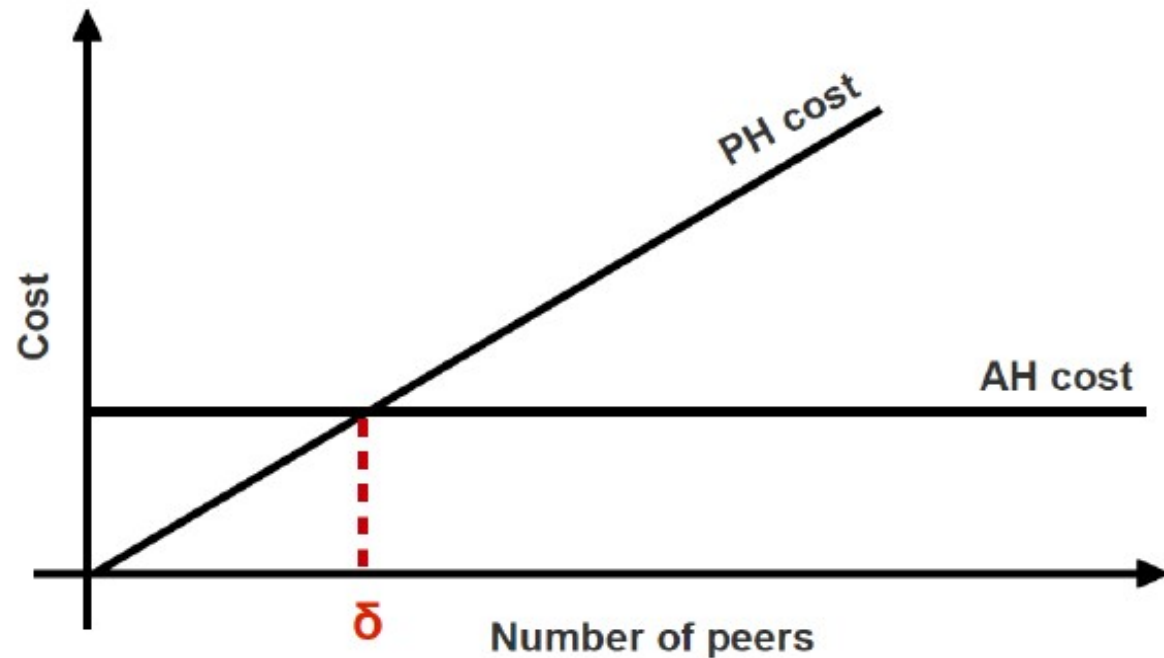
PH Cost

- PH cost in one round: $C_{ph} = C_{storage} + r \cdot (C_{chunk} + C_{req})$
 - $C_{storage}$: the storage cost
 - C_{chunk} : chunk transfer cost
 - C_{req} : chunk request
 - r : the number of retrieved chunks from PH in one round

δ

- We define δ as the number of peers that is economically reasonable to serve with PH utilization instead to run an additional AH for them.

- $\delta = C_{ah} / C_{ph}$



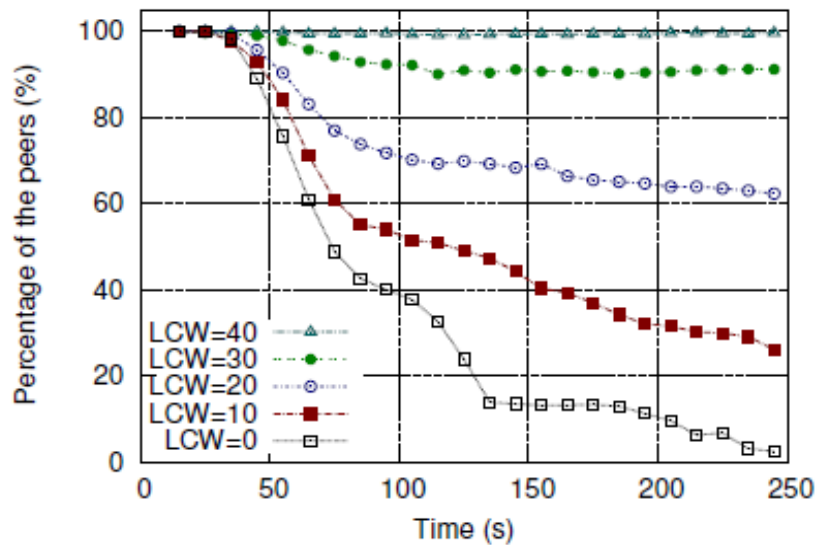
Manage the Cloud Resources

- If $\text{load} > \vartheta$: add AH
- If $\text{load} < \vartheta - H$: remove AH
 - H : number of peers served by one AH.
- Otherwise don't change AHs.

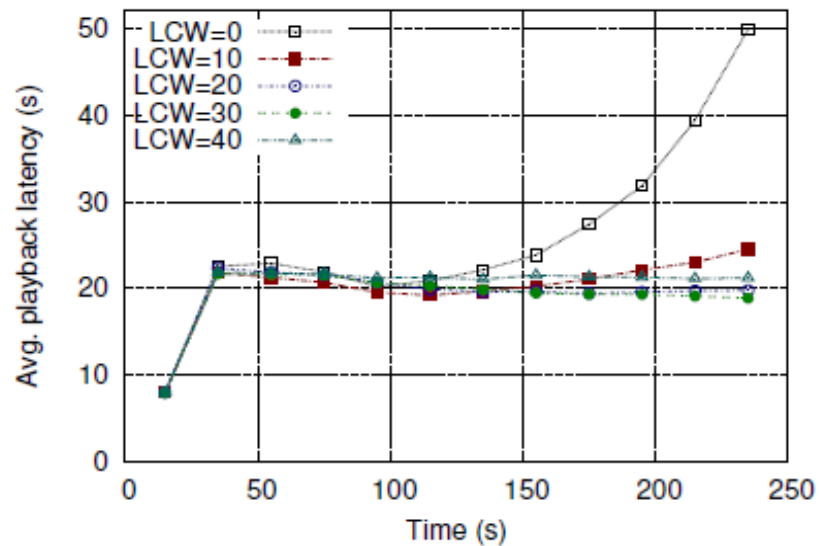
Experiments

Different LCW size

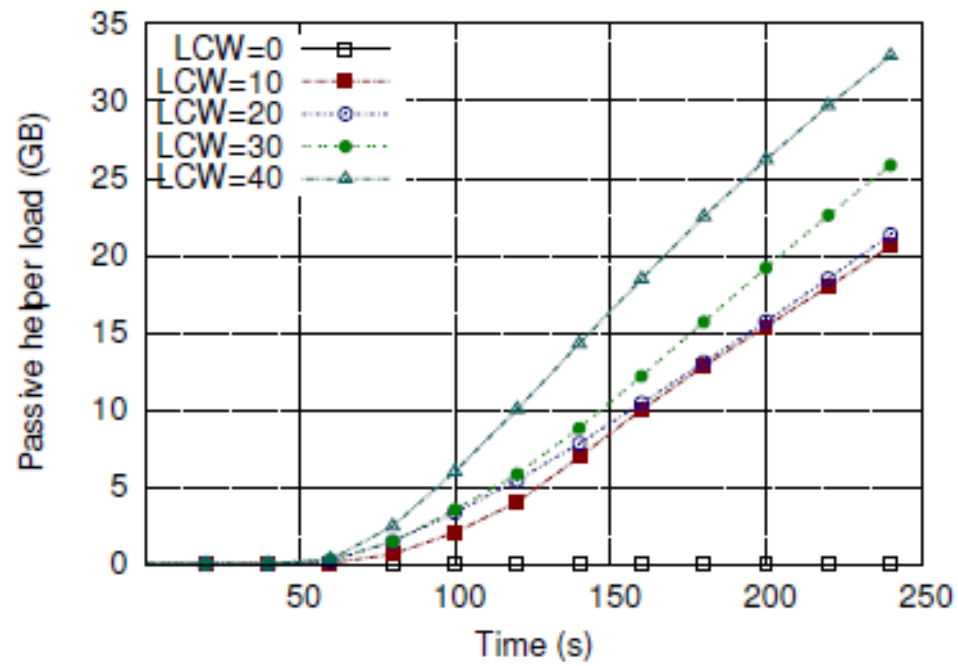
Playback continuity



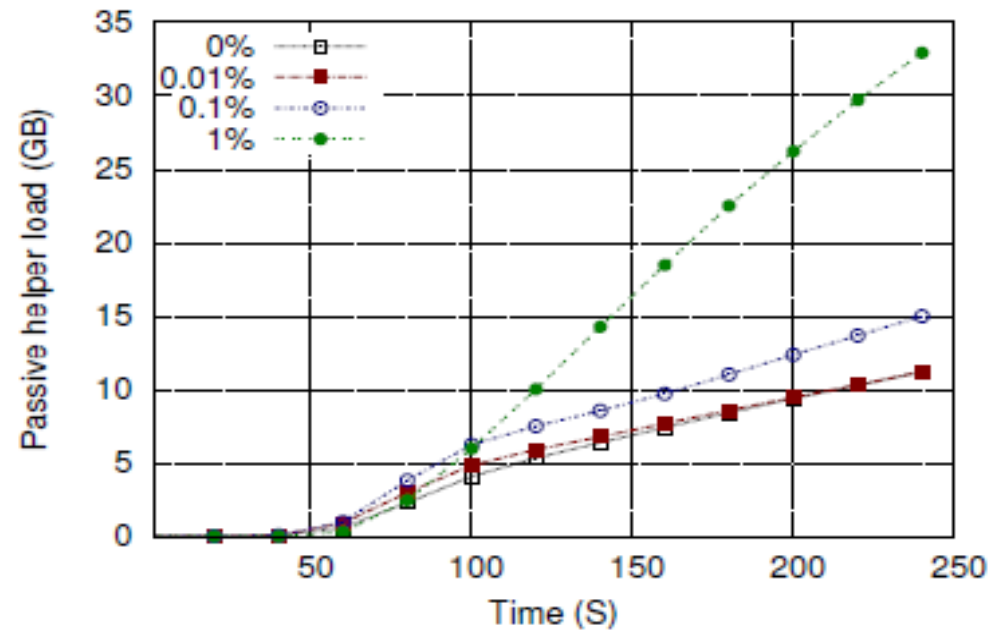
Playback latency



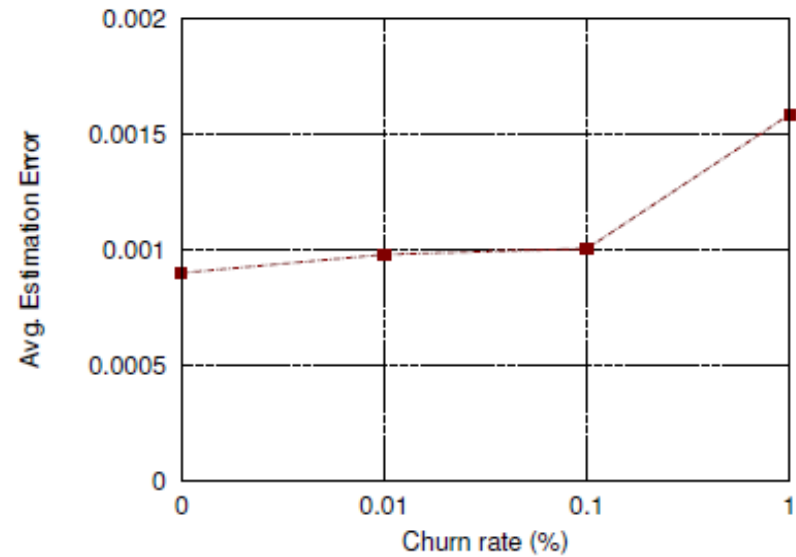
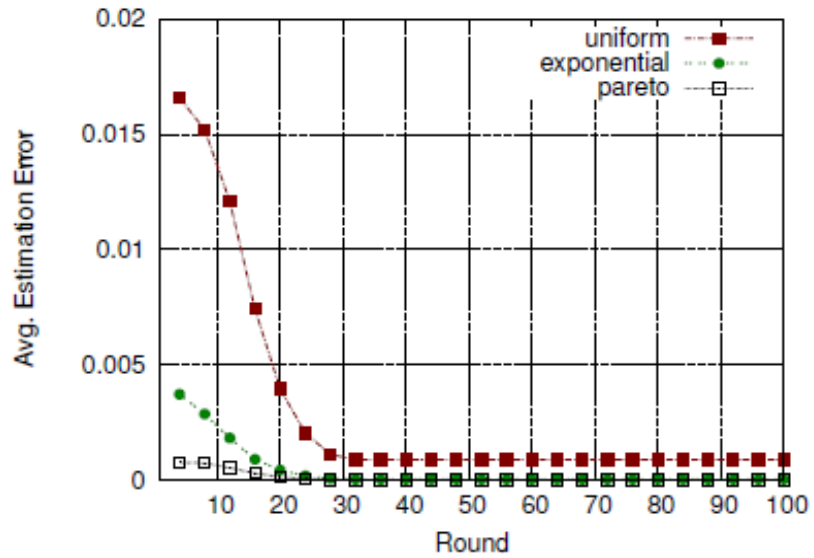
PH Load



PH Load in Different Churn Rates

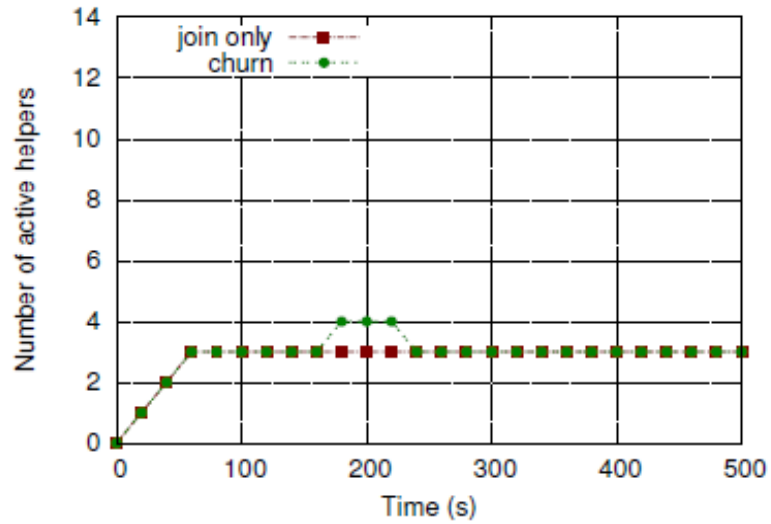


Slot Distribution Estimation

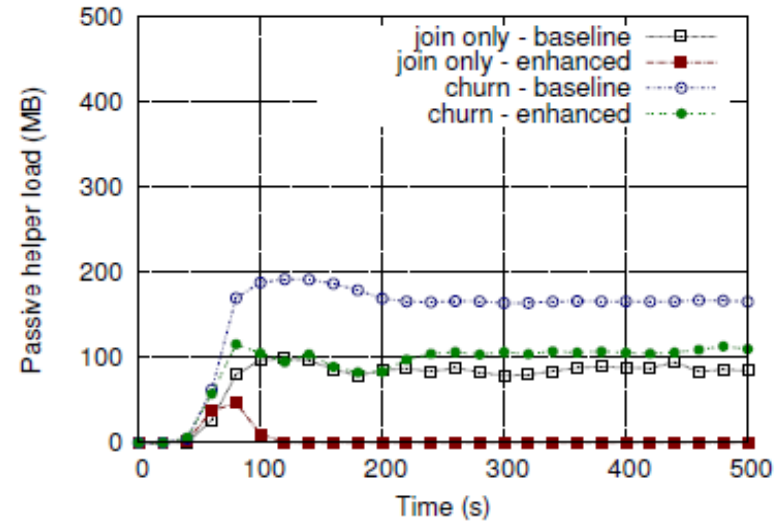


Cost

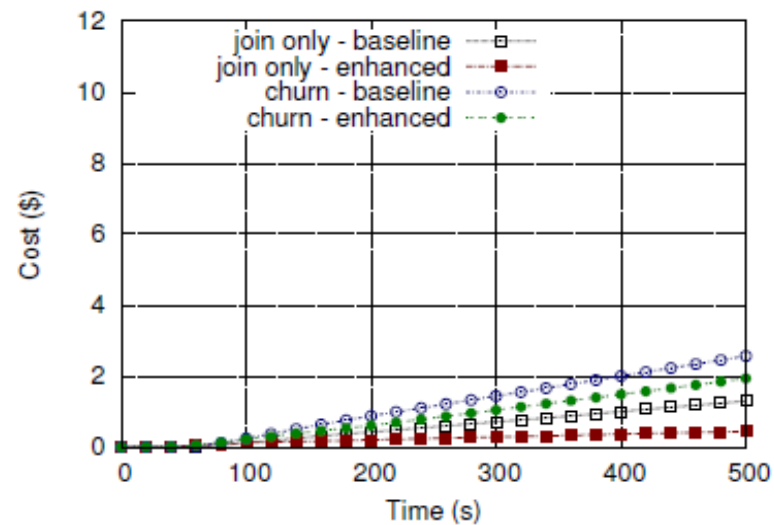
Number of AHs



PH load

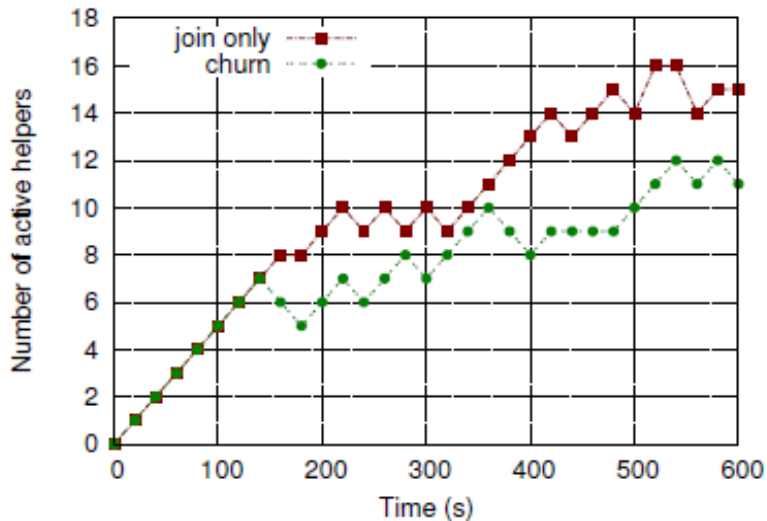


Total cost

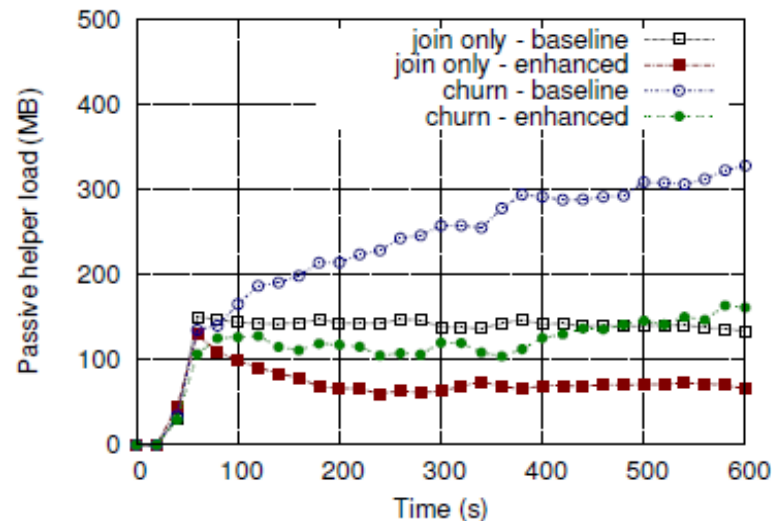


Cost – Real Trace Slot Distribution

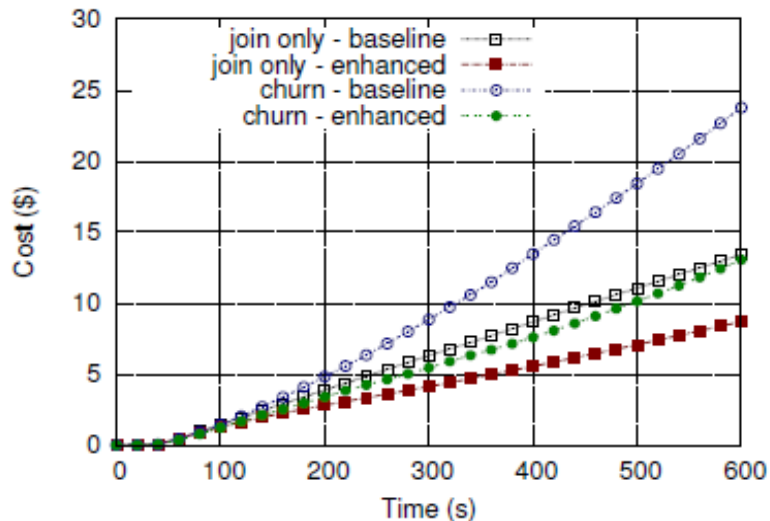
Number of AHs



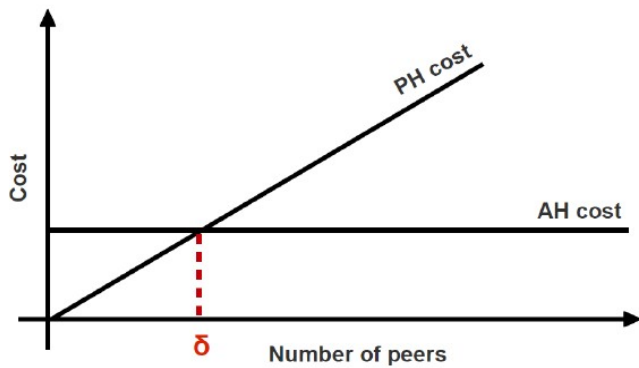
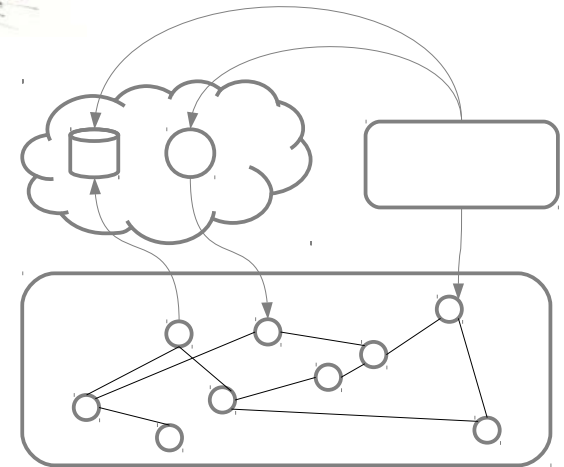
PH load



Total cost



Summary



Any Questions?